

# Sieves Talk - Aditya Ghosh

$p_n$  :  $n^{\text{th}}$  prime.

Twin prime conjecture :  $\liminf_{n \rightarrow \infty} (p_{n+1} - p_n) = 2$

GPY : (Goldston - Pintz - Yıldırım) (2005)

$$\liminf_{n \rightarrow \infty} \frac{p_{n+1} - p_n}{\log p_n} = 0$$

$$\left\{ \begin{array}{l} p_n \sim n \log n \\ \text{So Avg gap.} \approx \log n \end{array} \right.$$

Zhang : (2013) :  $\liminf_{n \rightarrow \infty} (p_{n+1} - p_n) \leq 7 \times 10^7$

Polymath :  $\leq 4600$

Maynard (2013) :  $\leq 600$

Polymath :  $\leq 246$

## Admissible ~~sets~~

$\{e.g.: \{0, 2\}$  not admissible,  $p: p+3 \pmod{2}$ )

DEF:  $H = \{h_1, \dots, h_k\}$  admissible if

$\forall$  primes  $p$ ,  $\exists n_p \in \mathbb{N}$  st  $n_p \not\equiv h_i \pmod{p}$   
 $\forall h_i$

[Prime  $k$ -tuples conjecture]:  $H$  be admissible.

Then  $\exists$  infinitely many  $n$  st  $\{n+h_1, \dots, n+h_k\}$  <sup>all</sup> primes

More generally, you can ask for linear forms  
 $a_i n + h_i$  or polynomials in  $n$ .  
 That basically gives you Goldbach's conjecture  
 & a lot of conj. in this field.

You can see easily this  $\left\{ \begin{array}{l} \Rightarrow \text{Twin primes conj (as } 0, 2 \text{ admissible)} \\ \Rightarrow \liminf_n (p_{n+m} - p_n) \leq (1+o(1)) m \log m \end{array} \right.$

|| G.P.Y.'s ~~method~~ : Studying approximations to this ||  
 || method ||  $k$ -tuple conjecture ||

## Primes in AP (Bombieri - Vinogradov thm)

$$(a, q) = 1,$$

$$\pi(x; q, a) = \#\{p \leq x : p \equiv a \pmod{q}\} \approx \frac{\pi(x)}{\phi(q)}$$

$$E_q := \sup_{a \text{ st } (a, q) = 1} \left| \pi(x; q, a) - \frac{\pi(x)}{\phi(q)} \right|$$

(error for each residue class)

DEF: (Level of distribution) Primes have l.o.d.  $\theta$  if

$$\forall A > 0,$$

$$\sum_{q \leq x^\theta} E_q \ll_A \frac{x}{(\log x)^A}$$

Thm: (Bombieri - Vinogradov) The primes have level of distribution

$$\theta \quad \forall \theta < \frac{1}{2}$$

{ "RH true on average" - avg over mod  $q$ .

|| ~~Elliot-Halberstam conj~~ Elliot-Halberstam conj: true for  $\forall \theta < 1$  ||

$\theta$  large: Strong level of equidistribution.

$\theta$  small: Small " ...

## Overview of the Sieve

①

|| Primes in AP ||  
 || (Bomb-Vin thm) ||

(Arithmetic result)

② Sieve Method (GPV sieve)

(Non-Arithmetic)

(i) Optimization problem: Choice of weight  
 (similar to Selberg sieve)

(ii) Combinatorial problem: Dense admissible sets  $H$

Zhang: Improves step ①

Maynard: (2)(i) More efficient sieve.

Multi-dimensional parameters for weight  
 (Since we have multiple elements in  $H$ )

{ Very flexible, can be generalized to other subsets  
 of primes, doesn't depend on  $E(H/O)$  as much }



LPT Sieve

$H = \{h_1, \dots, h_k\}$  admissible.

$$S = \sum_{n \leq x} \# \{i \mid n+h_i \text{ prime}\} w_n$$

$$\sum_{n \leq x} w_n$$

$w_n$ : Choice of weights  $> 0$

Analogous to what we did with the Selberg sieve

Think of it as a probability measure

$\times$  we're looking at expected no. of primes in this sequence.

Key obs: (i) If  $S > m$ , then  $\exists$  term  $n$  contributes  $> m$

(ii) Thus  $\exists$  at least  $m+1$  of the  $\{n+h_1, \dots, n+h_k\}$  which are primes.

(iii) If  $S > m$  for all large  $m$ , then

$$\liminf (p_{m+1} - p_m) \leq h_k - h_1 < \infty$$

We need  $S > 1$  for bounded gaps.

~~(Case  $e$  hypothesis)~~

$$\left[ \begin{array}{l} \text{Lem: } (h_1, \dots, h_k) \rightarrow k\text{-tuple } n \text{ in } [0, x] \text{ admissible} \\ \text{Then } (1+o(1)) \frac{x}{\log x} \leq k \leq \frac{x}{\log x} (2+o(1)) \\ h_k - h_1 \leq \frac{x}{\log x} > (\frac{1}{2} + o(1)) k \log k \end{array} \right]$$

Choice of  $w_n$ ~~no def~~

Manally,

(a)  $w_n$  small when few among the  $n+h_i$  are primes

(b) " large " a lot of " " "

(c) Find good upper bound for  $\sum_{N < n \leq 2N} w(n)$ good lower bound for  $\sum_{N < n \leq 2N} w(n) \#\{i \mid n+h_i \text{ prime}\}$ Guess 1:  $\prod_{j=1}^k \mathbb{1}_{n+h_j \text{ prime}}$ Had to find lower / upper bds  $i < c$ 

Guess 2: Sieves! (Selberg-like sieve)

~~Selberg sieve:~~  $w_n = \left( \sum_{\substack{d \mid (n+h_1) \dots (n+h_k) \\ d < R}} \lambda_d \right)^2$

(weights highly concentrated on primes.)

~~no def~~ $R$  small, so we care about small prime factors. $w_n$  small when  $(n+h_1) \dots (n+h_k)$  has many small prime factors.

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Choice of  $\lambda_d$  : Mimic Selberg's weight,

(i) Standard choice (inspired by other problems):

$$\lambda_d = \mu(d) \underbrace{\log(R/d)^k}_{\text{cutoff}}$$

If  $k$  large, then  $S \approx 0$   
 $\times EH(\theta)$

(primes have level 0)

So we just fail to prove bdd gaps  
 with  $\theta = 1 - \epsilon$

(Most optimistic estimate)

(ii) GPY choice: Optimize cutoff function  $f$ ,  $\lambda_d = \mu(d) f(d)$   
 Then,  $S \approx 2\theta$ .

So unconditionally ( $\theta = 1/2 - \epsilon$ ), just failed.

~~GPY~~ Zhang : Made  $EH(\theta)$  for  $\theta$  just a bit bigger than  $1/2$   
 $\times$  he could prove bdd gaps!!

[ARITHMETIC RESULT]

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Maynard's work

Generalizes the sieve with  $w_n$  to higher dimensions. parameters.

$$w_n = \sum_{\substack{d_1, \dots, d_k \\ d_i | n+1 \\ \prod d_i < \infty}} \lambda_{d_1, \dots, d_k} \approx \mu\left(\prod_{i=1}^k d_i\right) f(d_1, \dots, d_k)$$

This helps us look at divisions of  $n+1$  separately.  
Essentially, an optimization problem!  
(Use heuristics & calc of variations)

Given  $H = \{h_1, \dots, h_k\}$  admissible

$M_k = \sup_F \frac{J(F)}{I(F)}$   $J, I$  <sup>mimic</sup> the sums from before  
Integrals.

If  $M_k > 2m/\theta$ , then  $\exists$  infinitely many ~~prime~~ integers  $n$  s.t.  
at least  $m+1$  of  $n+h_i$  are primes.

Finding  $F$ .

$$F(t_1, \dots, t_k) = \begin{cases} \prod_{i=1}^k g(kt_i) & (t_1, \dots, t_k) \in \mathbb{Z}_{>0}^{(support)} \\ 0 & \text{else} \end{cases}$$

function  $g$

$k$  dim  $\rightarrow$  1 dim optimization problem.

~~$M_k$~~   $\rightarrow$  ~~log~~



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Q. In summary,

(i)  $M_k > \log k - 2 \log \log k - 2$  if  $k$  large

(ii) If  $M_k > 2m/\theta$ ,  $\exists$  int. many primes st at least  $m/\theta$  of the  $n$ th<sup>m</sup> primes.

lem: (i)  $\exists$  admissible set of size  $k$  contained in  $[0, H]$ ,  $H \approx k \log k$   
(take first  $k$  primes  $p > k$ )

(ii) We can take  $\theta < 1/2$  by Bomb vi.

$$\left\| \text{Thm: } \liminf_n (P_{n+m} - P_n) \leq cm^3 e^{4m} \right\|$$

Small k

$$F = \begin{cases} P(t_1, \dots, t_k) & (t_1, \dots, t_k) \in R_{1c} \\ 0 & \text{else} \end{cases}$$

$P$  symmetric poly  
(Stone-Weierstrass)

We can compute the integrals easily.

Prop: (i)  $M_{105} > 4$  ( $k=105$ ) ( $m=1$ )

(ii)  $\exists$  admissible set of size 105 in  $[0, 600]$  <sup>"H"</sup> (Computation)

Thus  $\liminf_{n \rightarrow \infty} P_{n+1} - P_n \leq 600$ .

Application  
~~of~~ of Maynard's work ~~and~~

$M_k \rightarrow \infty$  as  $k \rightarrow \infty$ , so ~~we~~ we have this  
 result for any  $\theta > 0$ .  
 Makes it very flexible. Don't need  $\theta = 1/2 - \epsilon$ .  
 So we can look at other subsets of primes, with weaker dist.

Consequences:

(i)  $d_n = p_{n+1} - p_n$ . There  $\exists$  arbitrarily long strings of  
 increasing gaps  $d_n < d_{n+1} < \dots < d_{n+m}$  in primes.

Similarly for decreasing

(ii)  $\exists$  arbitrarily large sets of primes with any pair differing  
 in at most 2 decimal places.